

Recall,

$$\gcd(a, b) = \max\{d : d|a \text{ and } d|b\}$$

Ex: $\gcd(22, 88)$

$$d|22 \Rightarrow d = 1, 2, 11, 22$$

$$d|88 \Rightarrow d = 1, 2, 4, 8, 11, 22, 44, 88$$

gcd Algorithm v.0:

goal: To find $\gcd(a, b)$

① Apply the fundamental theorem of arith.

⚠ This requires factoring which is HARD.

$$a = p_1^{e_1} p_2^{e_2} \dots p_n^{e_n} \text{ for } e_i \geq 0$$

$$b = p_1^{f_1} p_2^{f_2} \dots p_n^{f_n} \text{ for } f_j \geq 0$$

② Output $d = p_1^{\min\{e_1, f_1\}} p_2^{\min\{e_2, f_2\}} \dots p_n^{\min\{e_n, f_n\}}$

Ex: $\gcd(22, 88)$

$$= \gcd(2^1 \cdot 11^1, 2^3 \cdot 11^1)$$

$$= 2^{\min\{1, 3\}} \cdot 11^{\min\{1, 1\}}$$

$$= 2^1 \cdot 11^1 = 22$$

Ex: $\gcd(6, 35)$

$$= \gcd(2^1 3^1 5^0 7^0, 2^0 3^0 5^1 7^1)$$

$$= 2^{\min\{1, 0\}} 3^{\min\{1, 0\}} 5^{\min\{0, 1\}} 7^{\min\{0, 1\}}$$

$$= 2^0 \cdot 3^0 \cdot 5^0 \cdot 7^0 = 1$$

Lemma: $\gcd(a, b) = \gcd(a, b-a)$

we relate divisors of $(a \text{ and } b)$ to $(a \text{ and } b-a)$

If $d|a$ and $d|b$ then $d|(b-a)$

Thus, $\gcd(a, b) \leq \gcd(a, b-a)$

$S \subseteq T$
 $\Rightarrow \max S \leq \max T$

If $d|a$ and $d|(b-a)$ then $d|b$

Thus, $\gcd(a, b-a) \leq \gcd(a, b)$

GCD Algorithm v.1:

GOAL: To find $\gcd(a, b)$

If $a = b$ output a

If $a < b$ output $\gcd(a, b-a)$

If $a > b$ output $\gcd(a-b, a)$

⚠ RECURSION.

Ex: $\gcd(22, 88)$
 $= \gcd(22, 66)$
 $= \gcd(22, 44)$
 $= \gcd(22, 22)$
 $= 22$

Ex: $\gcd(13, 8)$
 $= \gcd(5, 8)$
 $= \gcd(5, 3)$
 $= \gcd(2, 3)$
 $= \gcd(2, 1)$
 $= \gcd(1, 1)$
 $= 1$

"Fib" Fact: If $a \leq b$ and $F_n < b \leq F_{n+1}$

Lamé's Thm then it takes at most n steps

COMPUTATIONAL COMPLEXITY to calculate $\gcd(a, b)$

Recall,

If $d|a$ and $d|b$ and $b = qa + r$
 then $d|r$

GCD Algorithm v.2:

GOAL: To find $\text{gcd}(a, b)$ where $a \leq b$ If $a = 0$ then output b .If $a = b$ then output a .If $a < b$ then

$$b = qa + r$$

and output $\text{gcd}(r, a)$ NB: $\text{gcd}(0, b) = b$ since $b|0$.Ex: Use the Euclidean Algorithm to compute $\text{gcd}(10, 103)$

$$103 = 10 \cdot 10 + 3 \quad \# \text{gcd}(103, 10)$$

$$10 = 3 \cdot 3 + 1 \quad \# \text{gcd}(3, 10)$$

$$3 = 3 \cdot 1 + 0 \quad \# \text{gcd}(3, 1)$$

Ex: Use the Euclidean algorithm to compute $\text{gcd}(13, 8)$

$$13 = 1 \cdot 8 + 5 \quad \# \text{gcd}(13, 8)$$

$$8 = 1 \cdot 5 + 3 \quad \# \text{gcd}(5, 8)$$

$$5 = 1 \cdot 3 + 2 \quad \# \text{gcd}(5, 3)$$

$$3 = 1 \cdot 2 + 1 \quad \# \text{gcd}(2, 3)$$

$$2 = 2 \cdot 1 + 0 \quad \# \text{gcd}(2, 1)$$

$$\# \text{gcd}(0, 1)$$

Defⁿ: A LINEAR DIOPHANTINE EQUATION has the form $Ax + By = C$ where $A, B, C \in \mathbb{Z}$ and we seek $x, y \in \mathbb{Z}$.

Ex: Use the Euc. Algo. to solve $13x + 5y = 1$.

apply the Euc. Algo. # Reverse the steps.

$$13 = 2 \cdot 5 + 3$$

$$\text{Thus, } 1 = 3 - 2$$

$$5 = 1 \cdot 3 + 2$$

$$= (13 - 2 \cdot 5) - (5 - 3)$$

$$= (13 - 2 \cdot 5) - (5 - [13 - 2 \cdot 5])$$

$$3 = 1 \cdot 2 + 1$$

$$= 2 \cdot 13 - 5 \cdot 5$$

Thus, $1 = 13 \cdot (2) + 5 \cdot (-5)$ is a solution.

Thm: $Ax + By = C$ has a solution iff $\gcd(A, B) \mid C$.

\Rightarrow] If $Ax + By = C$ has a solution then

$$d \mid A \text{ and } d \mid B \Rightarrow d \mid C$$

Thus, $\gcd(A, B) \mid C$

\Leftarrow] Use the Euclidean Algorithm to solve

$$Ax_1 + By_1 = \gcd(A, B)$$

Write $C = k \gcd(A, B)$ and obtain

$$A(kx_1) + B(ky_1) = C.$$

Ex: Use the Euc. Algo. to solve $107x + 10y = 1$.

$$107 = 10 \cdot 10 + 7$$

$$10 = 1 \cdot 7 + 3$$

$$7 = 2 \cdot 3 + 1$$

Thus,

$$1 = 7 - 2 \cdot 3$$

$$= (107 - 10 \cdot 10)$$

$$- 2(10 - 7)$$

$$= (107 - 10 \cdot 10)$$

$$- 2(10 - [107 - 10 \cdot 10])$$

$$= 3 \cdot 107 - 12 \cdot 10$$

Ex: Use the Euc. Algo. to solve $6x + 25y = 7$.

$$25 = 4 \cdot 6 + 1$$

$$6 = 6 \cdot 1 + 0$$

solve $1 = 6x + 25y$

$$1 = 25 - 4 \cdot 6$$

multiply through by 7.

$$7 = 7 \cdot 25 + 6 \cdot (-28)$$

Thus, $\gcd(25, 6) = 1$.

Summary:

— The Euclidean Algorithm (Practice!)

— Solving linear Diophantine equations

Ex: Write out the multiplication table mod 4.

•	1	2	3
1	1	2	3
2	2	0	2
3	3	2	1

NB: There is no x s.t.
 $2x \equiv 1 \pmod{4}$.

Ex: Write out the mult. table mod 5.

•	1	2	3	4
1	1	2	3	4
2	2	4	1	3
3	3	1	4	2
4	4	3	2	1

NB: There IS x s.t.
 $2x \equiv 1 \pmod{5}$.

Defⁿ: A **MULTIPLICATIVE INVERSE** of $k \pmod{n}$ is x such that $kx \equiv 1 \pmod{n}$.

"What is $\frac{1}{3}$?" $\frac{1}{3}$ is x such that $3x = 1$.

In this sense, $\frac{1}{3} \equiv 2 \pmod{5}$ since $2 \cdot 3 \equiv 1 \pmod{5}$.

Thm (Inverses are unique)

If $kx \equiv ky \equiv 1 \pmod{n}$ then $x \equiv y \pmod{n}$.

Pf: Suppose $kx \equiv 1 \pmod{n}$ and $ky \equiv 1 \pmod{n}$.

$$ky \equiv 1 \pmod{n} \Rightarrow xky \equiv x \cdot 1 \pmod{n}$$

$$\# \text{ using } xk \equiv 1 \pmod{n} \Rightarrow (xk)y \equiv x \pmod{n}$$

$$\Rightarrow y \equiv x \pmod{n}.$$

Thm: If $\gcd(a, m) = 1$ then $ax \equiv 1 \pmod{m}$ has a solution.

Pf: Use the Euclidean algorithm to solve

$$ax + my = 1.$$

$$1 \equiv ax + my \equiv ax \pmod{m}.$$

Thus x is a solution of $ax \equiv 1 \pmod{m}$.

Ex: Find the inverse of 6 mod 35.

$$35 = 5 \cdot 6 + 5$$

$$6 = 1 \cdot 5 + 1$$

$$\begin{aligned} \text{Thus, } 1 &= 6 - 5 \\ &= 6(35 - 5 \cdot 6) \end{aligned}$$

$$= 6 \cdot 6 - 35$$

$$\text{Thus, } 6 \cdot 6 \equiv 1 \pmod{35}.$$

Ex: Find the inverse of 5 mod 14

$$14 = 2 \cdot 5 + 4$$

$$5 = 1 \cdot 4 + 1$$

$$\text{Thus, } 1 = 5 - 4$$

$$= 5 - (14 - 2 \cdot 5)$$

$$= 3 \cdot 5 - 14.$$

$$\text{Thus } 3 \cdot 5 \equiv 1 \pmod{14}$$

Arithmetic Mod a Prime

Thm (Fermat): If $p \nmid a$ then $a^{p-1} \equiv 1 \pmod{p}$

NB: $p \nmid a \Rightarrow \gcd(a, p) = 1 \Rightarrow aA \equiv 1 \pmod{p}$

Pf: Consider the set $S = \{1, 2, 3, \dots, p-1\}$ and the map $x \mapsto ax \pmod{p}$ from S to S .

If $ax \equiv ay \pmod{p}$ then $Aax \equiv Aay \pmod{p}$
 $\Rightarrow x \equiv y \pmod{p}$.

Since $x, y \leq p-1$ we get $x = y$.

Thus, every element of S gets mapped to a distinct element of S . The map just re-arranges the set S .

Thus,

$$\begin{aligned} 1 \cdot 2 \cdot 3 \cdots (p-1) &\equiv (1 \cdot a)(2 \cdot a) \cdots [(p-1) \cdot a] \\ &\equiv [1 \cdot 2 \cdot 3 \cdots (p-1)] \cdot a^{p-1} \end{aligned}$$

cancel $1 \cdot 2 \cdot 3 \cdots (p-1)$ from both sides.

Therefore, $1 \equiv a^{p-1} \pmod{p}$.

Memorize this proof.

Thm (Wilson) p is prime $\implies (p-1)! \equiv p-1 \pmod{p}$.

Pf: If p is prime then $k \leq p-1 \implies \gcd(k, p) = 1$.

Thus, each term in $(p-1)!$ has an inverse.

These cancel in pairs, except $k=1$ and $k=p-1$.

Therefore, $(p-1)! \equiv (p-1) \pmod{p}$.

E.g.: $10! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 \cdot 6 \cdot 7 \cdot 8 \cdot 9 \cdot 10$

The diagram shows the product 10! = 1 · 2 · 3 · 4 · 5 · 6 · 7 · 8 · 9 · 10. Brackets are drawn under the terms to show pairings: (2, 6), (3, 4), (5, 9), and (7, 8). The term 10 is left unpaired.

$$= (2 \cdot 6) \cdot (3 \cdot 4) \cdot (5 \cdot 9) \cdot (7 \cdot 8) \cdot 10$$

$$\equiv 1 \cdot 1 \cdot 1 \cdot 1 \cdot 10 \equiv 11-1 \pmod{11}$$

Fact: $(4-1)! \equiv 2 \pmod{4}$.

Thm: If $n > 4$ is composite then $(n-1)! \equiv 0 \pmod{n}$.

If $n = pq$ where $p < q$ then $(n-1)! \equiv 1 \cdot 2 \cdot p \cdots q \cdots (n-1) \equiv 0 \pmod{n}$

If $n = p^2$ then $1 < p < 2p \leq n-1$ thus
 $(n-1)! = 1 \cdot 2 \cdot 3 \cdots p \cdots (2p) \cdots (n-1) \equiv 0 \pmod{n}$.

Silly Primality Test: compute $(k-1)! \pmod{k}$.

If you get 2, then $k=4$. If you get 0, then k is composite otherwise, k is prime.